

# A Deontic Logic of Knowingly Complying

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## Abstract

We introduce a logic for representing the deontic notion of *knowingly complying* –associated to an agent’s consciousness of taking a normative course of action for achieving a certain goal. Our logic features operators:

- for describing normative courses of actions, and
- for describing what each agent knowingly complies with.

We provide:

- a sound and complete axiom system,
- the computational complexity of its satisfiability problem, and
- an extension with an additional operator for capturing abilities (with a sound and complete axiom system).

## A Motivating Example

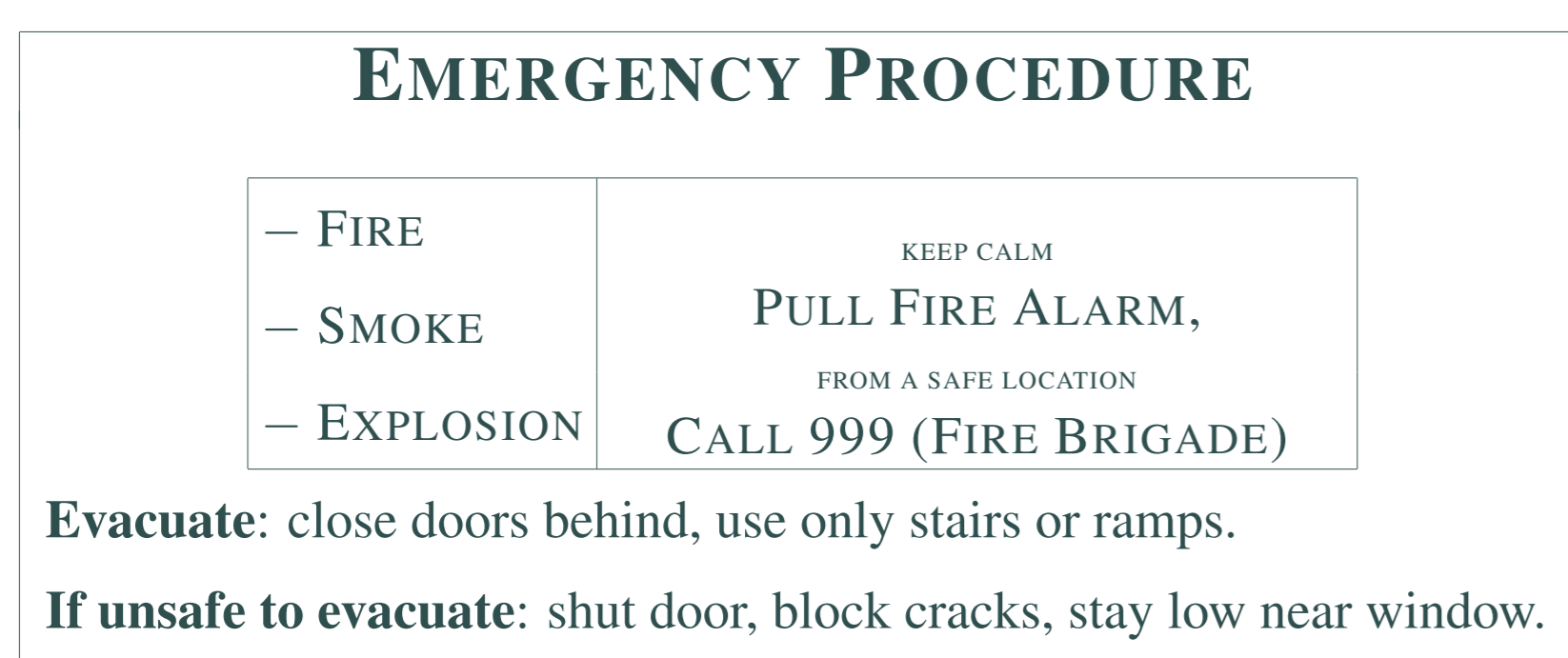


Figure 1: Fire Emergency Evacuation Plan

The main norms are:

- In the event of a fire/smoke/explosion, sound the nearest fire alarm, move to a safe location, and call the Fire Brigade.
- When evacuating close doors behind, use only stairs or ramps (never the elevator)
- Remain calm in any possible situation.

## The DLKc logic

**Definition 0.1.** Let  $\text{Prop}$  be a set for proposition symbols, and  $\text{Agt}$  a non-empty finite set of agent names. The language of DLKc is:

$$\varphi := p \mid \neg\varphi \mid \varphi \vee \psi \mid N(\psi, \varphi) \mid Kc_i(\psi, \varphi),$$

where:  $p \in \text{Prop}$  and  $i \in \text{Agt}$ .  $A\varphi = N(\neg\varphi, \perp)$  and  $E\varphi = \neg A\neg\varphi$ . Intuitively,  $N(\psi, \varphi)$ : “there is a normative course of action that brings about  $\varphi$  given  $\psi$ ”;  $Kc_i(\psi, \varphi)$ : “agent  $i$  knowingly complies with  $\varphi$  given  $\psi$ ”.

**Definition 0.2 (LTS).** A LTS is a tuple  $\mathcal{L} = \langle S, R, V \rangle$  where:

- $S$  is a non-empty set of states;
- $R = \{R_a \subseteq S^2 \mid a \in \text{Act}\}$ ; and
- $V : S \rightarrow 2^{\text{Prop}}$  is an assignment function.

**Definition 0.3 (Plans).** Let  $\text{Act}$  a set of basic actions,  $\pi \in \text{Act}^*$ . For  $0 \leq k \leq |\pi|$ ,  $\pi_k$  is the initial segment of  $\pi$  of length  $k$  and  $\pi[k]$  is the  $k^{\text{th}}$  element of  $\pi$ .

**Definition 0.4 (Strong Executability).** Let  $\mathcal{L} = \langle S, R, V \rangle$  be an LTS,  $s \in S$ ,  $\pi \in \text{Act}^*$ .  $\pi$  is SE at  $s$  iff for all  $k \in [0, |\pi| - 1]$  and all  $s' \in R_{\pi_k}(s)$ , implies  $R_{\pi[k+1]}(s') \neq \emptyset$ .  $\text{SE}(\pi) = \{s \in S \mid \pi \text{ is SE at } s\}$ .  $\Pi \subseteq \text{Act}^*$ ,  $\text{SE}(\Pi) = \bigcap_{\pi \in \Pi} \text{SE}(\pi)$ .

**Definition 0.5 (U-NLTS).** An uncertainty-based normative LTSs (U-NLTS) is a tuple  $\mathfrak{M} = \langle S, R, V, U, N \rangle$  where:

- $\mathcal{L} = \langle S, R, V \rangle$  is an LTS;
- $N \subseteq \text{Act}^*$  s.t. there is  $\pi \in N$  with  $\text{SE}(\pi) = S$ ;
- $U : \text{Agt} \rightarrow 2^{\text{Act}^*}$  satisfies:
  - $\emptyset \in U(i)$ , and
  - for all  $\{\Pi, \Pi'\} \subseteq U(i)$ ,  $\Pi \neq \Pi'$  implies  $\Pi \cap \Pi' = \emptyset$ .

**Definition 0.6 (Semantics).** Let  $\mathfrak{M} = \langle S, R, V, U, N \rangle$  be a U-NLTS,  $s \in S$ ,  $\varphi$  and  $\psi$  formulas:

$$\begin{aligned} \mathfrak{M}, s \Vdash p & \quad \text{iff } p \in V(s), \\ \mathfrak{M}, s \Vdash \neg\varphi & \quad \text{iff } \mathfrak{M}, s \not\Vdash \varphi, \\ \mathfrak{M}, s \Vdash \varphi \vee \psi & \quad \text{iff } \mathfrak{M}, s \Vdash \varphi \text{ or } \mathfrak{M}, s \Vdash \psi, \\ \mathfrak{M}, s \Vdash N(\psi, \varphi) & \quad \text{iff exists } \pi \in N \text{ such that} \end{aligned}$$

- (i)  $\llbracket \psi \rrbracket^{\mathfrak{M}} \subseteq \text{SE}(\pi)$  and
- (ii)  $R_{\pi}(\llbracket \psi \rrbracket^{\mathfrak{M}}) \subseteq \llbracket \varphi \rrbracket^{\mathfrak{M}}$ ,

$\mathfrak{M}, s \Vdash Kc_i(\psi, \varphi)$  iff exists  $\Pi \in U(i)$  such that

- (i)  $\Pi \subseteq N$ ,
- (ii)  $\llbracket \psi \rrbracket^{\mathfrak{M}} \subseteq \text{SE}(\Pi)$ , and
- (iii)  $R_{\Pi}(\llbracket \psi \rrbracket^{\mathfrak{M}}) \subseteq \llbracket \varphi \rrbracket^{\mathfrak{M}}$ ,

where  $\llbracket \chi \rrbracket^{\mathfrak{M}} = \{s \in S \mid \mathfrak{M}, s \Vdash \chi\}$ .

## A Motivating Example (cont.)

Fig. 1 can be represented as an U-NLTS  $\mathfrak{M} = \langle S, R, V, U, N \rangle$  s.t. the LTS  $\mathcal{L}$  part is modelled as in Fig. 2:

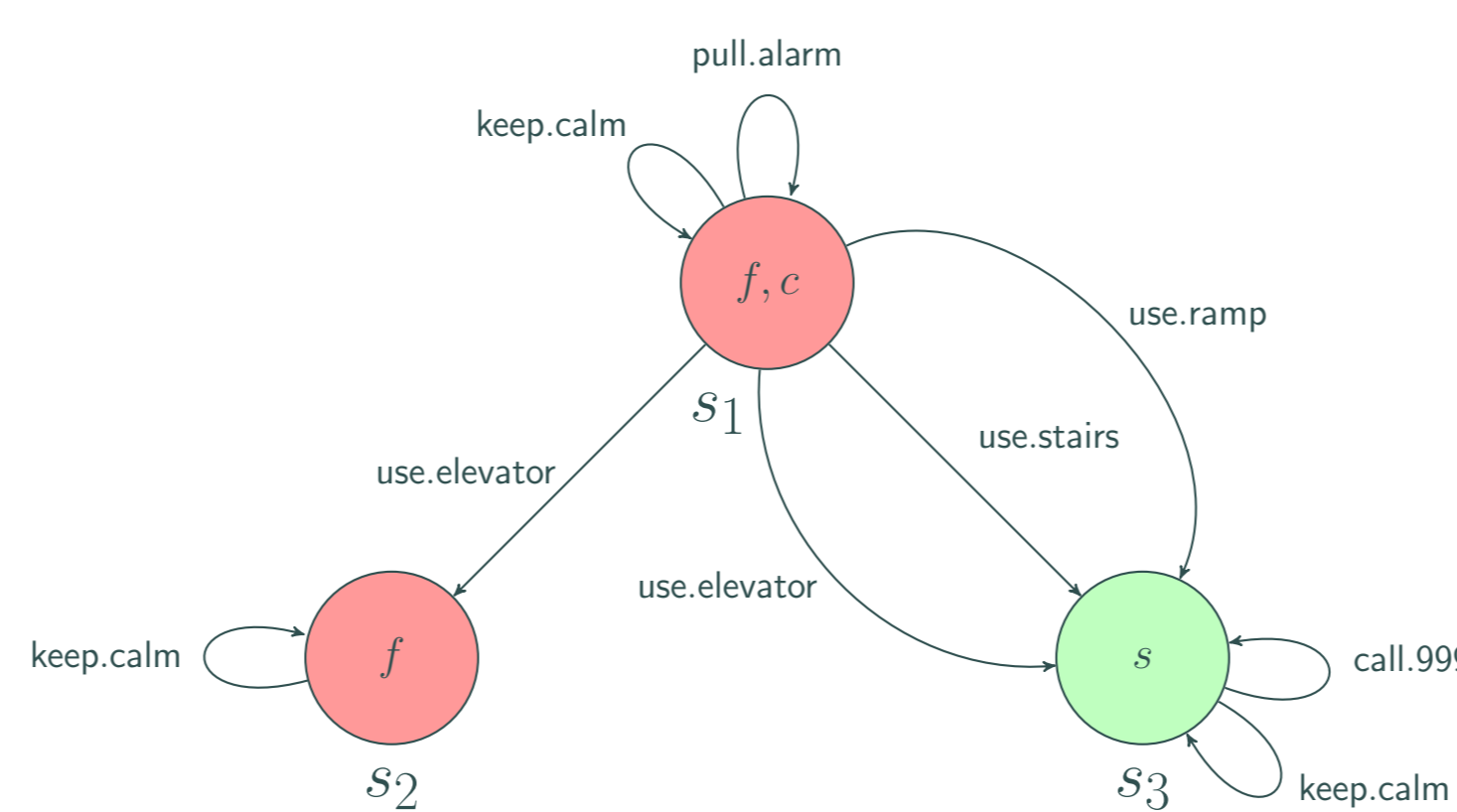


Figure 2: An LTS for the FEPP.

Where the basic actions are:

keep.calm, pull.alarm, call.999,  
use.stairs, use.ramp, use.elevator.

The plans considered are:

$\pi_0 = \text{keep.calm}$   
 $\pi_r = \text{pull.alarm; use.ramp; call.999}$   
 $\pi_s = \text{pull.alarm; use.stairs; call.999}$   
 $\pi_e = \text{pull.alarm; use.elevator; call.999}$

And each of the states in  $S$  represents a different situation:

- $s_1$ : a fire occurs ( $f$ ), there is the capacity to follow the FEPP ( $c$ ).
- $s_2$ : a fire occurs ( $f$ ), there is no capacity to follow the FEPP ( $\neg c$ ).
- $s_3$ : a safe location has been reached ( $s$ ), there is no fire ( $\neg f$ ).

Thus,  $\text{SE}(\pi_0) = S$ ,  $\text{SE}(\pi_r) = \text{SE}(\pi_s) = \{s_1\}$ , and  $\text{SE}(\pi_e) = \emptyset$  and the set of normative plans is  $N = \{\pi_0, \pi_r, \pi_s\}$ .

Suppose we have two agents:

- $i$ : has taken an occupational safety course and knows the difference between using stairs/ramps and using the elevator.

$$U(i) = \{\emptyset, \{\pi_s, \pi_r\}, \{\pi_e\}\}$$

- $j$ : has not taken the course and considers all possible ways of exiting the building might be equally good.

$$U(j) = \{\emptyset, \{\pi_s, \pi_e, \pi_r\}\}.$$

The following properties hold in  $\mathfrak{M}$

- (1)  $\mathfrak{M}, s_1 \Vdash A(s \rightarrow \neg f)$
- (2)  $\mathfrak{M}, s_1 \Vdash E f$
- (3)  $\mathfrak{M}, s_1 \Vdash N(f \wedge c, s)$
- (4)  $\mathfrak{M}, s_1 \Vdash Kc_i(f \wedge c, s)$
- (5)  $\mathfrak{M}, s_1 \not\Vdash Kc_j(f \wedge c, s)$

(1) and (2) are immediate. As a witness for (3) we can take the plan  $\pi_s$ . As a witness for (4) we can take the set  $\{\pi_s, \pi_r\} \in U(i)$ . Failure of (5) obtains from the fact that  $\{\pi_s, \pi_e, \pi_r\} \notin N$  and  $\llbracket f \wedge c \rrbracket^{\mathfrak{M}} \not\subseteq \text{SE}(\emptyset) = \emptyset$ .

## Axiomatization & complexity

Axioms:

Taut  $\vdash \varphi$  for  $\varphi$  a propositional tautology  
 DistA  $\vdash A(\psi \rightarrow \varphi) \rightarrow (A\psi \rightarrow A\varphi)$   
 TA  $\vdash A\varphi \rightarrow \varphi$

4KcA  $\vdash Kc_i(\psi, \varphi) \rightarrow AKc_i(\psi, \varphi)$   
 5KcA  $\vdash \neg Kc_i(\psi, \varphi) \rightarrow A\neg Kc_i(\psi, \varphi)$   
 4NA  $\vdash N(\psi, \varphi) \rightarrow AN(\psi, \varphi)$   
 5NA  $\vdash \neg N(\psi, \varphi) \rightarrow A\neg N(\psi, \varphi)$

KcN  $\vdash Kc_i(\psi, \varphi) \rightarrow N(\psi, \varphi)$

DN  $\vdash N(\varphi, \top)$

KcA  $\vdash (A(\psi \rightarrow \chi) \wedge Kc_i(\chi, \rho) \wedge A(\rho \rightarrow \varphi)) \rightarrow Kc_i(\psi, \varphi)$

NA  $\vdash (A(\psi \rightarrow \chi) \wedge N(\chi, \rho) \wedge A(\rho \rightarrow \varphi)) \rightarrow N(\psi, \varphi)$

Kc $\perp$   $\vdash Kc_i(\perp, \perp)$

Rules:

$$\frac{\vdash \psi \quad \vdash (\psi \rightarrow \varphi)}{\vdash \varphi} \text{ (MP)} \quad \frac{\vdash \varphi}{\vdash A\varphi} \text{ (Nec)}$$

Table 1: Axiom system  $\mathcal{DLKc}$  for DLKc over U-NLTSs.

**Theorem 1.** The axiom system  $\mathcal{DLKc}$  in Tab. 1 is sound and strongly complete for DLKc over the class of all U-NLTSs.

**Proposition 0.1.** The model checking problem for DLKc is in P.

**Theorem 2.** The satisfiability problem for DLKc is NP-complete.

## Reasoning About Abilities

**Definition 0.7.** The language of  $\text{DLKc}^+$  is defined by:

$$\varphi := p \mid \neg\varphi \mid \varphi \vee \psi \mid S(\psi, \varphi) \mid N(\psi, \varphi) \mid Kc_i(\psi, \varphi),$$

where:  $p \in \text{Prop}$  and  $i \in \text{Agt}$ . Intuitively,  $S(\psi, \varphi)$ : “there is a course of action that brings about  $\varphi$  given  $\psi$ ”.

**Definition 0.8.** Let  $\mathfrak{M} = \langle S, R, V, U, N \rangle$  be a U-NLTS,  $s \in S$ ,  $\psi$  and  $\varphi$  formulas:

$\mathfrak{M}, s \Vdash S(\psi, \varphi)$  iff exists  $\pi \in \text{Act}^*$  such that

- (i)  $\llbracket \psi \rrbracket^{\mathfrak{M}} \subseteq \text{SE}(\pi)$  and
- (ii)  $R_{\pi}(\llbracket \psi \rrbracket^{\mathfrak{M}}) \subseteq \llbracket \varphi \rrbracket^{\mathfrak{M}}$ .

4SA  $\vdash S(\psi, \varphi) \rightarrow AS(\psi, \varphi)$       NS  $\vdash N(\psi, \varphi) \rightarrow S(\psi, \varphi)$   
 5SA  $\vdash \neg S(\psi, \varphi) \rightarrow A\neg S(\psi, \varphi)$       EmpS  $\vdash A(\psi \rightarrow \varphi) \rightarrow S(\psi, \varphi)$   
 CompS  $\vdash (S(\psi, \chi) \wedge S(\chi, \varphi)) \rightarrow S(\psi, \varphi)$

Table 2: Additional axioms for  $\text{DLKc}^+$ .

**Theorem 3.** The axioms and rules in Tabs. 1 and 2 yield a sound and strongly complete axiom system for  $\text{DLKc}^+$  over the class of all U-NLTSs.

## Future work

- Characterize the exact complexity of the satisfiability problem of the extended logic  $\text{DLKc}^+$ .
- Establish different levels of responsibility for the agents using the relation between the set of plans  $U(i)$  of each agent and the set of norms  $N$ .
- Impose new restrictions on the different components of the model (or weakening them), and obtain new logics.

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